



# CeTA 3.0 – Improved Support for ARI, CoCo and Infeasibility

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joint work with Akihisa Yamada, Dohan Kim and Teppei Saito supported by the Austrian Science Fund (FWF) project I 5943

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  - commutation: ✓ CPP 2024 on parallel critical pairs (Dohan, Kiraku, Nao, René)
  - non-commutation: this talk, ongoing



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- B.4: formalize techniques for rewriting induction and ground confluence
  - sorted rewriting and pattern completeness (√?) FSCD 2024 (Akihisa, René)
  - soundness of rewriting induction (Akihisa, Dohan, René): ongoing work
  - generation of certificates (Naoki, Takahito): format not fixed

#### CeTA 2.x – Problems prior to ARI-Format

- consider commutation and GCR
  - common property: signature matters
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- resolving conflict in CeTA 2.x: maximally accepting
  - in commutation proofs, signature consists of symbols in input
  - in non-commutation proofs, signature is not restricted
- $\implies$  in CeTA 2.x it might have been possible to accept commutation proof and non-commutation proof of same two TRSs due to different implicit signature

# CeTA 3.0 – Improved Support for ARI

- ARI format explicitly contains signature
- improved support:
  - CeTA 3.0 reads signature from ARI format
  - all proofs in CeTA using ARI format are now taking care of proper handling of signature
    - COM
    - INF
    - CR: signature ignored, but CR has signature extension property
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  - complexity proofs always allowed signature (to define basic terms)
  - termination proofs now consistently use "signature is not restricted" semantics

- initial design of CPF used in CeTA 2.x
  - one self-contained CPF, containing input and proof, e.g.
    - TRS and non-termination proof
    - CTRS and confluence proof
  - similar to Isabelle theories, which also contain mixture of specs and proofs
  - no synchronization problem, as in "CR proof of COPS 120"

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is COPS 120 confluent?

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- competition management somehow needs to extract input from certificate and then find conflict "COPS 120  $\neq$  {}"
- Akihisa's idea: instead of extracting input, let CeTA take input separately from competition management software

- CeTA 3.0 extends Akihisa's idea, and splits a CPF into four parts
  - input: a TRS, two TRSs, CTRS + infeasibility query, ...
  - property: termination, confluence, ground-confluence, commutation, ...
  - answer: yes, no, upperbound  $\mathcal{O}(n^2), \ldots$
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  - property is CR, but CPF contains a termination proof
  - input was some TRS, but CPF contains a proof for different TRS (if the same proof can be used for both TRSs, then this is accepted)

- CPF 2 had several inconsistencies or non-uniform treatments
  - sometimes removed rules had to be specified, sometimes remaining rules
  - four different formats to specify joinable critical pairs, e.g. in <ruleLabeling>,
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     <parallelClosed>, <pcpClosed>, <decreasingDiagrams>
- in CPF 3 and CeTA 3.0 the format has been simplified and unified
  - always specify removed rules (decreases certificate size from  $\mathcal{O}(n^2)$  to  $\mathcal{O}(n)$ )
  - uniform way to specify joining sequences, choose between
    - *left*,  $t_1$ ,  $t_2$ , ...,  $t_n$ , *right* intermediate terms suffice
    - specify upper bounds on steps bfs; fails on conversions that are not joins
    - for WCR only: "rewrite to normal form"

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  - optional specification of a term index and a rule index
    - example: in compositional confluence criteria, one can specify

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  - CPFs of termCOMP 2023:
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• CeTA 3.0 directly expands parsed indices (future work)

#### CPF 2.0 vs 3.0

- demo
  - Example 39 of FSCD 2022 paper (Kiraku, Nao)
  - some rule-labeling proof (by Julian Nagele)



# CeTA 3.0 – Improved Support for Infeasibility (and Non-CR and Non-COM)

- common theme for all these properties: show non-reachability property
  - non-CR and non-COM: given peak  $s \xleftarrow{*}{\mathcal{P}} u \xrightarrow{*}{s} t$ , show that

$$s \xrightarrow{*}{\mathcal{S}} v \xleftarrow{*}{\mathcal{R}} t$$

is impossible

• given oriented infeasibility query  $s_1 \approx t_1, \ldots, s_n \approx t_n$ , define  $s := c(s_1, \ldots, s_n)$ and  $t := c(t_1, \ldots, t_n)$  for fresh symbol c and show that

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• available semantic solutions (Akihisa, Takahito): find discrimination pair (non-CR, non-COM) or co-rewrite pair (infeasibility) and solve some constraints involving  $\mathcal{R}$ ,  $\mathcal{S}$ , s, t, e.g.,  $\mathcal{R} \subseteq \succeq$  and  $t \succ s$  for infeasibility

# Discrimination Pairs and Co-Rewrite Pairs

- co-rewrite pair ( $\succ$ ,  $\succeq$ )
  - $\succeq \cap \prec = \emptyset$
  - $\succeq$  is transitive and reflexive, closed under contexts and substitutions
  - > is irreflexive and closed under substitutions
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- CeTA 3.0
  - complete redesign of interface for relations on terms
  - basic properties can individually be demanded
  - wrapper functions for common cases

#### New Interface

- considers three relations (S = strict, NS = non-strict, NST = non-strict top)
- simplified(!) properties (dropped argument filters, Ce-compatibility,...)

```
"rel impl prop impl \equiv \exists S NS NST.
     — <implementation approximates real relations>
      (\forall \text{ st. (isOK(rel impl.s impl st)} \longrightarrow \text{ st} \in S) \land
               (isOK(rel impl.ns impl st) \longrightarrow st \in NS) \land
               (isOK(rel impl.nst impl st) \longrightarrow st \in NST))
     — (unconditional properties)
     \land irrefl S
     ∧ ctxt.closed NS
     ∧ subst closed NS
     \wedge trans NS
     ∧ refl NS
     — <properties that can be tested via flags>
     \land (isOK(rel impl.standard impl) \rightarrow
              trans S \land S \subseteq NS \land S 0 NS \subseteq S \land NS 0 S \subseteq S \land subst.closed NST
           \land trans NST \land NST 0 S \subset S \land S 0 NST \subset S)
     \land (isOK(rel impl.mono impl []) \rightarrow ctxt.closed S)
     \land (isOK(rel impl.top mono impl) \longrightarrow top mono NS NST)
     \land (isOK(rel impl.top refl impl) \rightarrow refl NST)
      \land (isOK(rel impl.SN impl) \longrightarrow SN S)
     \land (isOK(rel impl.subst s impl) \rightarrow subst.closed S)
     \land (isOK(rel impl.co rewr impl) \longrightarrow NS \cap S<sup>-1</sup> = {})
      \land (\forall cm cc. isOK(rel impl.cpx impl cm cc) \rightarrow deriv bound measure class S cm cc)"
```

### Co-Rewrite Pairs in New Interface

```
    specification via interface is simple

definition rel impl co rewrite pair where
 "rel impl co rewrite pair impl = do {
      rel impl.co rewr impl;
      rel impl.subst s impl
    } <+? (\lambda . showsl lit (STR ''error message ''))"
```

```
    property is easy to use
```

innsbruck

```
lemma rel impl co rewrite pair: assumes "rel impl prop impl"
  and "isOK(rel impl co rewrite pair impl)"
  and "isOK(rel impl s impl s constraints)" "isOK(rel impl ns impl ns constraints)"
shows "\exists S NS. co rewrite pair S NS \land set s constraints \subseteq S \land set ns constraints \subseteq NS"
```

#### discrimination pairs are similar

- WPO has been generalized in formalization (René)
  - example: SN of underlying order is propagated, but not demanded
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you are invited to add these relations to your certificate generating tool; increase the power of certified non-CR, non-COM, infeasibility proving



# Summary

- big restructuring efforts have gone into CeTA 3.0 and CPF 3.0
- achieved
  - better support for certification in competitions
  - more consistency in CPF 3
  - reduced size of certificates
  - new term orders became available
- $\implies$  ask for support of CPF 3 in tools
- $\implies$  ask for support of these in tools



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- unclear
  - plans to publish restructuring efforts
    - is there a plan to write ARI-infrastructure paper?
    - $\implies$  interest: add section on new certification approach
  - plans to publish formalization of new orders
    - let's discuss among Akihisa, Dohan, René, Teppei

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Questions?



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